

Indexing Highly Dynamic Hierarchical Data

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Challenge: Complex Updates on Hierarchical Data

Leaf

relocate_leaf (E, below J)

Subtree

relocate_subtree (C, below J)

Inner Node

relocate_inner (C, [F, J])

Range

relocate_range ([B, F], below J)

Asymptotic Runtime

Scheme	Update Operations				
	leaf	subtree	inner node	range	skew(u)
Adjacency	1	1	c	c	1
Linked	1	1	c	c	1
NI	n	n	n	n	n
Dyn-NI	1	s	1	s	u
Dyn-NI-Parent	1	s	c	s	u
Dyn-NI-Level	1	s	c	s	u
Dewey path	l	l(f+s)	l(f+s)	l(f+s)	l(f+s)
Dyn-Dewey	l	ls	ls	ls	l+u
AO-Tree	1	log n	log n	log n	1
BO-Tree[B]	1	B log _B n	B log _B n	B log _B n	1
O-List[B]	1	s/B+B	s/B+B	s/B+B	u/B ²
DeltaNI	log u	log u	log u	log u	log u

■ efficient
 ■ mostly efficient
 ■ inefficient

Problem: Labeling schemes cannot handle complex updates efficiently!

Relocate C below H

Key	Parent	PSL	NI	GapNI	Ordpath
A	null	0;8;0	[0,17]	[0,1700]	1
B	A	1;0;1	[1,2]	[100,200]	1.1
C	A	2;4;1	[3,12]	[300,1200]	1.3
D	C	3;0;2	[4,5]	[400,500]	1.3.1
E	C	4;0;2	[6,7]	[600,700]	1.3.3
F	C	5;0;2	[8,9]	[800,900]	1.3.5
G	C	6;0;2	[10,11]	[1000,1100]	1.3.7
H	A	7;0;1	[13,14]	[1300,1400]	1.5
I	A	8;0;1	[15,16]	[1500,1600]	1.7

Solution: Order Indexes

Concept: Encode nested intervals by positions in an ordered data structure

Hierarchy

Table

RID	Key	level	lower	upper
0	A	0	0	17
1	B	1	1	2
2	C	2	3	12
3	D	3	4	5
4	E	3	6	7
5	F	3	8	9
6	G	3	10	11

Order Index

Efficient Updates

Delete subtree (B)

Split

Join

Flavors: AO-Tree, BO-Tree, Order List

AO-Tree

BO-Tree

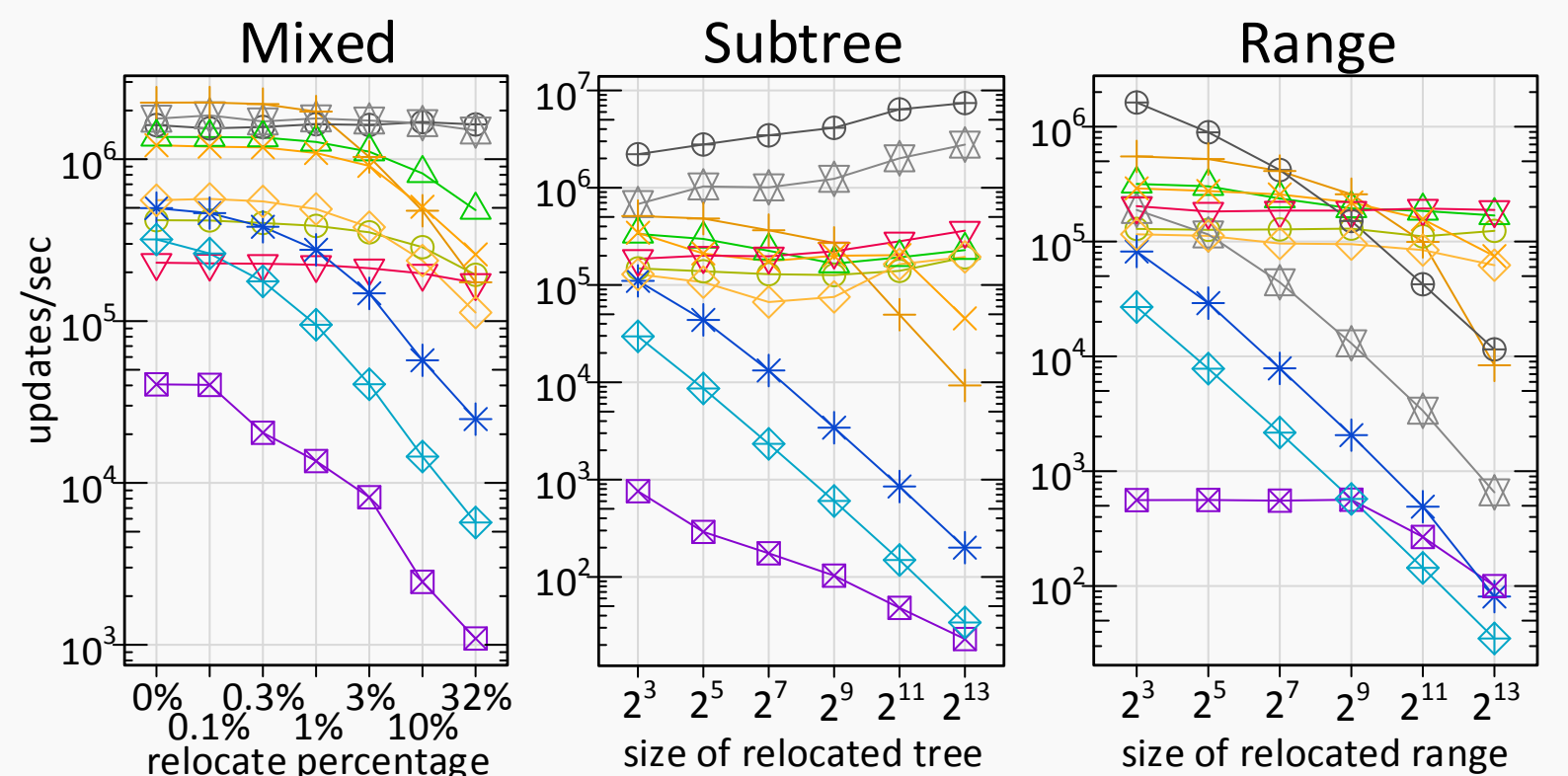
Order List

Queries

← level([4]) ← before([0,3]) ← before([3,10])

Performance

Updates: Complex updates facilitated efficiently



Queries: On a par with labeling schemes

